

Matrix

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theory MatrixGeneral imports Main begin

types 'a infmatrix = [nat, nat]  $\Rightarrow$  'a

constdefs
  nonzero-positions :: ('a::zero) infmatrix  $\Rightarrow$  (nat*nat) set
  nonzero-positions A == {pos. A (fst pos) (snd pos)  $\sim$  0}

typedef 'a matrix = {(f::('a::zero) infmatrix)). finite (nonzero-positions f)}
  <proof>

declare Rep-matrix-inverse[simp]

lemma finite-nonzero-positions : finite (nonzero-positions (Rep-matrix A))
  <proof>

constdefs
  nrows :: ('a::zero) matrix  $\Rightarrow$  nat
  nrows A == if nonzero-positions(Rep-matrix A) = {} then 0 else Suc(Max
    ((image fst) (nonzero-positions (Rep-matrix A))))
  ncols :: ('a::zero) matrix  $\Rightarrow$  nat
  ncols A == if nonzero-positions(Rep-matrix A) = {} then 0 else Suc(Max ((image
    snd) (nonzero-positions (Rep-matrix A))))

lemma nrows:
  assumes hyp: nrows A  $\leq$  m
  shows (Rep-matrix A m n) = 0 (is ?concl)
  <proof>

constdefs
  transpose-infmatrix :: 'a infmatrix  $\Rightarrow$  'a infmatrix
  transpose-infmatrix A j i == A i j
  transpose-matrix :: ('a::zero) matrix  $\Rightarrow$  'a matrix
  transpose-matrix == Abs-matrix o transpose-infmatrix o Rep-matrix

declare transpose-infmatrix-def[simp]
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lemma *transpose-infmatrix-twice[simp]*: *transpose-infmatrix (transpose-infmatrix A) = A*
 <proof>

lemma *transpose-infmatrix*: *transpose-infmatrix (% j i. P j i) = (% j i. P i j)*
 <proof>

lemma *transpose-infmatrix-closed[simp]*: *Rep-matrix (Abs-matrix (transpose-infmatrix (Rep-matrix x))) = transpose-infmatrix (Rep-matrix x)*
 <proof>

lemma *infmatrixforward*: *(x::'a infmatrix) = y \implies \forall a b. x a b = y a b* <proof>

lemma *transpose-infmatrix-inject*: *(transpose-infmatrix A = transpose-infmatrix B) = (A = B)*
 <proof>

lemma *transpose-matrix-inject*: *(transpose-matrix A = transpose-matrix B) = (A = B)*
 <proof>

lemma *transpose-matrix[simp]*: *Rep-matrix(transpose-matrix A) j i = Rep-matrix A i j*
 <proof>

lemma *transpose-transpose-id[simp]*: *transpose-matrix (transpose-matrix A) = A*
 <proof>

lemma *nrows-transpose[simp]*: *nrows (transpose-matrix A) = ncols A*
 <proof>

lemma *ncols-transpose[simp]*: *ncols (transpose-matrix A) = nrows A*
 <proof>

lemma *ncols*: *ncols A <= n \implies Rep-matrix A m n = 0*
 <proof>

lemma *ncols-le*: *(ncols A <= n) = (! j i. n <= i \longrightarrow (Rep-matrix A j i) = 0) (is - = ?st)*
 <proof>

lemma *less-ncols*: *(n < ncols A) = (? j i. n <= i & (Rep-matrix A j i) \neq 0) (is ?concl)*
 <proof>

lemma *le-ncols*: *(n <= ncols A) = (\forall m. (\forall j i. m <= i \longrightarrow (Rep-matrix A j i) = 0) \longrightarrow n <= m) (is ?concl)*
 <proof>

lemma *nrows-le*: $(\text{nrows } A \leq n) = (! j \ i. n \leq j \longrightarrow (\text{Rep-matrix } A \ j \ i) = 0)$
(is ?s)
 <proof>

lemma *less-nrows*: $(m < \text{nrows } A) = (? j \ i. m \leq j \ \& \ (\text{Rep-matrix } A \ j \ i) \neq 0)$
(is ?concl)
 <proof>

lemma *le-nrows*: $(n \leq \text{nrows } A) = (\forall m. (\forall j \ i. m \leq j \longrightarrow (\text{Rep-matrix } A \ j \ i) = 0) \longrightarrow n \leq m)$ **(is ?concl)**
 <proof>

lemma *nrows-notzero*: $\text{Rep-matrix } A \ m \ n \neq 0 \implies m < \text{nrows } A$
 <proof>

lemma *ncols-notzero*: $\text{Rep-matrix } A \ m \ n \neq 0 \implies n < \text{ncols } A$
 <proof>

lemma *finite-natarray1*: $\text{finite } \{x. x < (n::\text{nat})\}$
 <proof>

lemma *finite-natarray2*: $\text{finite } \{\text{pos}. (\text{fst pos}) < (m::\text{nat}) \ \& \ (\text{snd pos}) < (n::\text{nat})\}$
 <proof>

lemma *RepAbs-matrix*:
assumes *aem*: $? m. ! j \ i. m \leq j \longrightarrow x \ j \ i = 0$ **(is ?em)** **and** *aen*: $? n. ! j \ i. (n \leq i \longrightarrow x \ j \ i = 0)$ **(is ?en)**
shows $(\text{Rep-matrix } (\text{Abs-matrix } x)) = x$
 <proof>

constdefs

apply-infmatrix :: $('a \Rightarrow 'b) \Rightarrow 'a \text{ infmatrix} \Rightarrow 'b \text{ infmatrix}$
apply-infmatrix *f* == $\% A. (\% j \ i. f \ (A \ j \ i))$
apply-matrix :: $('a \Rightarrow 'b) \Rightarrow ('a::\text{zero}) \text{ matrix} \Rightarrow ('b::\text{zero}) \text{ matrix}$
apply-matrix *f* == $\% A. \text{Abs-matrix } (\text{apply-infmatrix } f \ (\text{Rep-matrix } A))$
combine-infmatrix :: $('a \Rightarrow 'b \Rightarrow 'c) \Rightarrow 'a \text{ infmatrix} \Rightarrow 'b \text{ infmatrix} \Rightarrow 'c \text{ infmatrix}$
combine-infmatrix *f* == $\% A \ B. (\% j \ i. f \ (A \ j \ i) \ (B \ j \ i))$
combine-matrix :: $('a \Rightarrow 'b \Rightarrow 'c) \Rightarrow ('a::\text{zero}) \text{ matrix} \Rightarrow ('b::\text{zero}) \text{ matrix} \Rightarrow ('c::\text{zero}) \text{ matrix}$
combine-matrix *f* == $\% A \ B. \text{Abs-matrix } (\text{combine-infmatrix } f \ (\text{Rep-matrix } A) \ (\text{Rep-matrix } B))$

lemma *expand-apply-infmatrix[simp]*: $\text{apply-infmatrix } f \ A \ j \ i = f \ (A \ j \ i)$
 <proof>

lemma *expand-combine-infmatrix[simp]*: $\text{combine-infmatrix } f \ A \ B \ j \ i = f \ (A \ j \ i) \ (B \ j \ i)$
 <proof>

constdefs

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commutative :: ('a ⇒ 'a ⇒ 'b) ⇒ bool
commutative f == ! x y. f x y = f y x
associative :: ('a ⇒ 'a ⇒ 'a) ⇒ bool
associative f == ! x y z. f (f x y) z = f x (f y z)

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To reason about associativity and commutativity of operations on matrices, let's take a step back and look at the general situation: Assume that we have sets A and B with $B \subset A$ and an abstraction $u : A \rightarrow B$. This abstraction has to fulfill $u(b) = b$ for all $b \in B$, but is arbitrary otherwise. Each function $f : A \times A \rightarrow A$ now induces a function $f' : B \times B \rightarrow B$ by $f' = u \circ f$. It is obvious that commutativity of f implies commutativity of f' : $f'xy = u(fxy) = u(fyx) = f'yx$.

lemma *combine-infmatrix-commute*:

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commutative f ⇒ commutative (combine-infmatrix f)
⟨proof⟩

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lemma *combine-matrix-commute*:

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commutative f ⇒ commutative (combine-matrix f)
⟨proof⟩

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On the contrary, given an associative function f we cannot expect f' to be associative. A counterexample is given by $A = \mathbb{Z}$, $B = \{-1, 0, 1\}$, as f we take addition on \mathbb{Z} , which is clearly associative. The abstraction is given by $u(a) = 0$ for $a \notin B$. Then we have

$$f'(f'11) - 1 = u(f(u(f11)) - 1) = u(f(u2) - 1) = u(f0 - 1) = -1,$$

but on the other hand we have

$$f'1(f'1 - 1) = u(f1(u(f1 - 1))) = u(f10) = 1.$$

A way out of this problem is to assume that $f(A \times A) \subset A$ holds, and this is what we are going to do:

lemma *nonzero-positions-combine-infmatrix[simp]*: $f\ 0\ 0 = 0 \Rightarrow \text{nonzero-positions } (\text{combine-infmatrix } f\ A\ B) \subseteq (\text{nonzero-positions } A) \cup (\text{nonzero-positions } B)$
 ⟨proof⟩

lemma *finite-nonzero-positions-Rep[simp]*: $\text{finite } (\text{nonzero-positions } (\text{Rep-matrix } A))$
 ⟨proof⟩

lemma *combine-infmatrix-closed [simp]*:

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f 0 0 = 0 ⇒ Rep-matrix (Abs-matrix (combine-infmatrix f (Rep-matrix A)
(Rep-matrix B))) = combine-infmatrix f (Rep-matrix A) (Rep-matrix B)
⟨proof⟩

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We need the next two lemmas only later, but it is analog to the above one, so we prove them now:

lemma *nonzero-positions-apply-infmatrix[simp]: $f \ 0 = 0 \implies \text{nonzero-positions } (\text{apply-infmatrix } f \ A) \subseteq \text{nonzero-positions } A$*
 ⟨proof⟩

lemma *apply-infmatrix-closed [simp]:*
 $f \ 0 = 0 \implies \text{Rep-matrix } (\text{Abs-matrix } (\text{apply-infmatrix } f \ (\text{Rep-matrix } A))) = \text{apply-infmatrix } f \ (\text{Rep-matrix } A)$
 ⟨proof⟩

lemma *combine-infmatrix-assoc[simp]: $f \ 0 \ 0 = 0 \implies \text{associative } f \implies \text{associative } (\text{combine-infmatrix } f)$*
 ⟨proof⟩

lemma *comb: $f = g \implies x = y \implies f \ x = g \ y$*
 ⟨proof⟩

lemma *combine-matrix-assoc: $f \ 0 \ 0 = 0 \implies \text{associative } f \implies \text{associative } (\text{combine-matrix } f)$*
 ⟨proof⟩

lemma *Rep-apply-matrix[simp]: $f \ 0 = 0 \implies \text{Rep-matrix } (\text{apply-matrix } f \ A) \ j \ i = f \ (\text{Rep-matrix } A \ j \ i)$*
 ⟨proof⟩

lemma *Rep-combine-matrix[simp]: $f \ 0 \ 0 = 0 \implies \text{Rep-matrix } (\text{combine-matrix } f \ A \ B) \ j \ i = f \ (\text{Rep-matrix } A \ j \ i) \ (\text{Rep-matrix } B \ j \ i)$*
 ⟨proof⟩

lemma *combine-nrows: $f \ 0 \ 0 = 0 \implies \text{nrows } (\text{combine-matrix } f \ A \ B) \leq \max (\text{nrows } A) (\text{nrows } B)$*
 ⟨proof⟩

lemma *combine-ncols: $f \ 0 \ 0 = 0 \implies \text{ncols } (\text{combine-matrix } f \ A \ B) \leq \max (\text{ncols } A) (\text{ncols } B)$*
 ⟨proof⟩

lemma *combine-nrows: $f \ 0 \ 0 = 0 \implies \text{nrows } A \leq q \implies \text{nrows } B \leq q \implies \text{nrows } (\text{combine-matrix } f \ A \ B) \leq q$*
 ⟨proof⟩

lemma *combine-ncols: $f \ 0 \ 0 = 0 \implies \text{ncols } A \leq q \implies \text{ncols } B \leq q \implies \text{ncols } (\text{combine-matrix } f \ A \ B) \leq q$*
 ⟨proof⟩

constdefs

zero-r-neutral :: ('a \Rightarrow 'b::zero \Rightarrow 'a) \Rightarrow bool
zero-r-neutral f == ! a. f a 0 = a

$zero-l-neutral :: ('a::zero \Rightarrow 'b \Rightarrow 'b) \Rightarrow bool$
 $zero-l-neutral\ f == !\ a.\ f\ 0\ a = a$
 $zero-closed :: (('a::zero) \Rightarrow ('b::zero) \Rightarrow ('c::zero)) \Rightarrow bool$
 $zero-closed\ f == (!x.\ f\ x\ 0 = 0) \ \&\ (!y.\ f\ 0\ y = 0)$

consts $foldseq :: ('a \Rightarrow 'a \Rightarrow 'a) \Rightarrow (nat \Rightarrow 'a) \Rightarrow nat \Rightarrow 'a$
primrec
 $foldseq\ f\ s\ 0 = s\ 0$
 $foldseq\ f\ s\ (Suc\ n) = f\ (s\ 0)\ (foldseq\ f\ (\% k.\ s(Suc\ k))\ n)$

consts $foldseq-transposed :: ('a \Rightarrow 'a \Rightarrow 'a) \Rightarrow (nat \Rightarrow 'a) \Rightarrow nat \Rightarrow 'a$
primrec
 $foldseq-transposed\ f\ s\ 0 = s\ 0$
 $foldseq-transposed\ f\ s\ (Suc\ n) = f\ (foldseq-transposed\ f\ s\ n)\ (s\ (Suc\ n))$

lemma $foldseq-assoc : associative\ f \Longrightarrow foldseq\ f = foldseq-transposed\ f$
 $\langle proof \rangle$

lemma $foldseq-distr: \llbracket associative\ f; commutative\ f \rrbracket \Longrightarrow foldseq\ f\ (\% k.\ f\ (u\ k)\ (v\ k))\ n = f\ (foldseq\ f\ u\ n)\ (foldseq\ f\ v\ n)$
 $\langle proof \rangle$

theorem $\llbracket associative\ f; associative\ g; \forall a\ b\ c\ d.\ g\ (f\ a\ b)\ (f\ c\ d) = f\ (g\ a\ c)\ (g\ b\ d); ?\ x\ y.\ (f\ x) \neq (f\ y); ?\ x\ y.\ (g\ x) \neq (g\ y); f\ x\ x = x; g\ x\ x = x \rrbracket \Longrightarrow f=g \mid (!\ y.\ f\ y\ x = y) \mid (!\ y.\ g\ y\ x = y)$
 $\langle proof \rangle$

lemma $foldseq-zero:$
assumes $fz: f\ 0\ 0 = 0$ **and** $sz: !\ i.\ i \leq n \longrightarrow s\ i = 0$
shows $foldseq\ f\ s\ n = 0$
 $\langle proof \rangle$

lemma $foldseq-significant-positions:$
assumes $p: !\ i.\ i \leq N \longrightarrow S\ i = T\ i$
shows $foldseq\ f\ S\ N = foldseq\ f\ T\ N$ (**is** $?concl$)
 $\langle proof \rangle$

lemma $foldseq-tail: M \leq N \Longrightarrow foldseq\ f\ S\ N = foldseq\ f\ (\% k.\ (if\ k < M\ then\ (S\ k)\ else\ (foldseq\ f\ (\% k.\ S(k+M))\ (N-M))))\ M$ (**is** $?p \Longrightarrow ?concl$)
 $\langle proof \rangle$

lemma $foldseq-zerotail:$
assumes
 $fz: f\ 0\ 0 = 0$
and $sz: !\ i.\ n \leq i \longrightarrow s\ i = 0$
and $nm: n \leq m$
shows
 $foldseq\ f\ s\ n = foldseq\ f\ s\ m$

$\langle \text{proof} \rangle$

lemma *foldseq-zerotail2*:

assumes $! x. f\ x\ 0 = x$

and $! i. n < i \longrightarrow s\ i = 0$

and $nm: n \leq m$

shows

$foldseq\ f\ s\ n = foldseq\ f\ s\ m$ (**is** $?concl$)

$\langle \text{proof} \rangle$

lemma *foldseq-zerostart*:

$! x. f\ 0\ (f\ 0\ x) = f\ 0\ x \implies ! i. i \leq n \longrightarrow s\ i = 0 \implies foldseq\ f\ s\ (Suc\ n) = f\ 0\ (s\ (Suc\ n))$

$\langle \text{proof} \rangle$

lemma *foldseq-zerostart2*:

$! x. f\ 0\ x = x \implies ! i. i < n \longrightarrow s\ i = 0 \implies foldseq\ f\ s\ n = s\ n$

$\langle \text{proof} \rangle$

lemma *foldseq-almostzero*:

assumes $f0x: ! x. f\ 0\ x = x$ **and** $fx0: ! x. f\ x\ 0 = x$ **and** $s0: ! i. i \neq j \longrightarrow s\ i = 0$

shows $foldseq\ f\ s\ n = (\text{if } (j \leq n) \text{ then } (s\ j) \text{ else } 0)$ (**is** $?concl$)

$\langle \text{proof} \rangle$

lemma *foldseq-distr-unary*:

assumes $!! a\ b. g\ (f\ a\ b) = f\ (g\ a)\ (g\ b)$

shows $g(foldseq\ f\ s\ n) = foldseq\ f\ (\% x. g(s\ x))\ n$ (**is** $?concl$)

$\langle \text{proof} \rangle$

constdefs

$mult_matrix_n :: nat \Rightarrow ((a::zero) \Rightarrow (b::zero) \Rightarrow (c::zero)) \Rightarrow ('c \Rightarrow 'c \Rightarrow 'c) \Rightarrow 'a\ matrix \Rightarrow 'b\ matrix \Rightarrow 'c\ matrix$

$mult_matrix_n\ n\ fmul\ fadd\ A\ B == Abs_matrix(\% j\ i. foldseq\ fadd\ (\% k. fmul\ (Rep_matrix\ A\ j\ k)\ (Rep_matrix\ B\ k\ i))\ n)$

$mult_matrix :: ((a::zero) \Rightarrow (b::zero) \Rightarrow (c::zero)) \Rightarrow ('c \Rightarrow 'c \Rightarrow 'c) \Rightarrow 'a\ matrix \Rightarrow 'b\ matrix \Rightarrow 'c\ matrix$

$mult_matrix\ fmul\ fadd\ A\ B == mult_matrix_n\ (\max\ (ncols\ A)\ (nrows\ B))\ fmul\ fadd\ A\ B$

lemma *mult-matrix-n*:

assumes $prems: ncols\ A \leq n$ (**is** $?An$) $nrows\ B \leq n$ (**is** $?Bn$) $fadd\ 0\ 0 = 0$ $fmul\ 0\ 0 = 0$

shows $c: mult_matrix\ fmul\ fadd\ A\ B = mult_matrix_n\ n\ fmul\ fadd\ A\ B$ (**is** $?concl$)

$\langle \text{proof} \rangle$

lemma *mult-matrix-nm*:

assumes $prems: ncols\ A \leq n$ $nrows\ B \leq n$ $ncols\ A \leq m$ $nrows\ B \leq m$ $fadd\ 0\ 0 = 0$ $fmul\ 0\ 0 = 0$

shows $mult_matrix_n\ n\ fmul\ fadd\ A\ B = mult_matrix_n\ m\ fmul\ fadd\ A\ B$

$\langle proof \rangle$

constdefs

r-distributive :: (*'a* \Rightarrow *'b* \Rightarrow *'b*) \Rightarrow (*'b* \Rightarrow *'b* \Rightarrow *'b*) \Rightarrow bool
r-distributive fmul fadd == ! *a u v*. *fmul a (fadd u v) = fadd (fmul a u) (fmul a v)*
l-distributive :: (*'a* \Rightarrow *'b* \Rightarrow *'a*) \Rightarrow (*'a* \Rightarrow *'a* \Rightarrow *'a*) \Rightarrow bool
l-distributive fmul fadd == ! *a u v*. *fmul (fadd u v) a = fadd (fmul u a) (fmul v a)*
distributive :: (*'a* \Rightarrow *'a* \Rightarrow *'a*) \Rightarrow (*'a* \Rightarrow *'a* \Rightarrow *'a*) \Rightarrow bool
distributive fmul fadd == *l-distributive fmul fadd* & *r-distributive fmul fadd*

lemma *max1*: !! *a x y*. (*a::nat*) $\leq x \Rightarrow a \leq \max x y$ $\langle proof \rangle$

lemma *max2*: !! *b x y*. (*b::nat*) $\leq y \Rightarrow b \leq \max x y$ $\langle proof \rangle$

lemma *r-distributive-matrix*:

assumes *prems*:

r-distributive fmul fadd

associative fadd

commutative fadd

fadd 0 0 = 0

! *a*. *fmul a 0 = 0*

! *a*. *fmul 0 a = 0*

shows *r-distributive (mult-matrix fmul fadd) (combine-matrix fadd)* (**is** ?*concl*)

$\langle proof \rangle$

lemma *l-distributive-matrix*:

assumes *prems*:

l-distributive fmul fadd

associative fadd

commutative fadd

fadd 0 0 = 0

! *a*. *fmul a 0 = 0*

! *a*. *fmul 0 a = 0*

shows *l-distributive (mult-matrix fmul fadd) (combine-matrix fadd)* (**is** ?*concl*)

$\langle proof \rangle$

instance *matrix* :: (*zero*) *zero* $\langle proof \rangle$

defs(**overloaded**)

zero-matrix-def: (*0::('a::zero) matrix*) == *Abs-matrix*(% *j i*. *0*)

lemma *Rep-zero-matrix-def[simp]*: *Rep-matrix 0 j i = 0*

$\langle proof \rangle$

lemma *zero-matrix-def-nrows[simp]*: *nrows 0 = 0*

$\langle proof \rangle$

lemma *zero-matrix-def-ncols[simp]*: *ncols 0 = 0*

$\langle \text{proof} \rangle$

lemma *combine-matrix-zero-l-neutral*: $\text{zero-l-neutral } f \implies \text{zero-l-neutral } (\text{combine-matrix } f)$
 $\langle \text{proof} \rangle$

lemma *combine-matrix-zero-r-neutral*: $\text{zero-r-neutral } f \implies \text{zero-r-neutral } (\text{combine-matrix } f)$
 $\langle \text{proof} \rangle$

lemma *mult-matrix-zero-closed*: $\llbracket \text{fadd } 0 \ 0 = 0; \text{zero-closed } \text{fmul} \rrbracket \implies \text{zero-closed}$
 $(\text{mult-matrix } \text{fmul } \text{fadd})$
 $\langle \text{proof} \rangle$

lemma *mult-matrix-n-zero-right[simp]*: $\llbracket \text{fadd } 0 \ 0 = 0; !a. \text{fmul } a \ 0 = 0 \rrbracket \implies$
 $\text{mult-matrix-n } n \text{ fmul fadd } A \ 0 = 0$
 $\langle \text{proof} \rangle$

lemma *mult-matrix-n-zero-left[simp]*: $\llbracket \text{fadd } 0 \ 0 = 0; !a. \text{fmul } 0 \ a = 0 \rrbracket \implies$
 $\text{mult-matrix-n } n \text{ fmul fadd } 0 \ A = 0$
 $\langle \text{proof} \rangle$

lemma *mult-matrix-zero-left[simp]*: $\llbracket \text{fadd } 0 \ 0 = 0; !a. \text{fmul } 0 \ a = 0 \rrbracket \implies \text{mult-matrix}$
 $\text{fmul fadd } 0 \ A = 0$
 $\langle \text{proof} \rangle$

lemma *mult-matrix-zero-right[simp]*: $\llbracket \text{fadd } 0 \ 0 = 0; !a. \text{fmul } a \ 0 = 0 \rrbracket \implies \text{mult-matrix}$
 $\text{fmul fadd } A \ 0 = 0$
 $\langle \text{proof} \rangle$

lemma *apply-matrix-zero[simp]*: $f \ 0 = 0 \implies \text{apply-matrix } f \ 0 = 0$
 $\langle \text{proof} \rangle$

lemma *combine-matrix-zero*: $f \ 0 \ 0 = 0 \implies \text{combine-matrix } f \ 0 \ 0 = 0$
 $\langle \text{proof} \rangle$

lemma *transpose-matrix-zero[simp]*: $\text{transpose-matrix } 0 = 0$
 $\langle \text{proof} \rangle$

lemma *apply-zero-matrix-def[simp]*: $\text{apply-matrix } (\% x. 0) \ A = 0$
 $\langle \text{proof} \rangle$

constdefs

singleton-matrix :: $\text{nat} \Rightarrow \text{nat} \Rightarrow ('a::\text{zero}) \Rightarrow 'a \text{ matrix}$
singleton-matrix $j \ i \ a == \text{Abs-matrix } (\% m \ n. \text{if } j = m \ \& \ i = n \text{ then } a \text{ else } 0)$
move-matrix :: $('a::\text{zero}) \text{ matrix} \Rightarrow \text{int} \Rightarrow \text{int} \Rightarrow 'a \text{ matrix}$
move-matrix $A \ y \ x == \text{Abs-matrix } (\% j \ i. \text{if } (\text{neg } ((\text{int } j) - y)) \mid (\text{neg } ((\text{int } i) - x))$
then $0 \text{ else Rep-matrix } A \ (\text{nat } ((\text{int } j) - y)) \ (\text{nat } ((\text{int } i) - x)))$
take-rows :: $('a::\text{zero}) \text{ matrix} \Rightarrow \text{nat} \Rightarrow 'a \text{ matrix}$

$\text{take-rows } A \ r == \text{Abs-matrix}(\% j \ i. \text{if } (j < r) \text{ then } (\text{Rep-matrix } A \ j \ i) \text{ else } 0)$
 $\text{take-columns} :: ('a::\text{zero}) \text{ matrix} \Rightarrow \text{nat} \Rightarrow 'a \text{ matrix}$
 $\text{take-columns } A \ c == \text{Abs-matrix}(\% j \ i. \text{if } (i < c) \text{ then } (\text{Rep-matrix } A \ j \ i) \text{ else } 0)$

constdefs

$\text{column-of-matrix} :: ('a::\text{zero}) \text{ matrix} \Rightarrow \text{nat} \Rightarrow 'a \text{ matrix}$
 $\text{column-of-matrix } A \ n == \text{take-columns } (\text{move-matrix } A \ 0 \ (- \text{int } n)) \ 1$
 $\text{row-of-matrix} :: ('a::\text{zero}) \text{ matrix} \Rightarrow \text{nat} \Rightarrow 'a \text{ matrix}$
 $\text{row-of-matrix } A \ m == \text{take-rows } (\text{move-matrix } A \ (- \text{int } m) \ 0) \ 1$

lemma $\text{Rep-singleton-matrix[simp]}: \text{Rep-matrix } (\text{singleton-matrix } j \ i \ e) \ m \ n = (\text{if } j = m \ \& \ i = n \text{ then } e \text{ else } 0)$
 $\langle \text{proof} \rangle$

lemma $\text{apply-singleton-matrix[simp]}: f \ 0 = 0 \implies \text{apply-matrix } f \ (\text{singleton-matrix } j \ i \ x) = (\text{singleton-matrix } j \ i \ (f \ x))$
 $\langle \text{proof} \rangle$

lemma $\text{singleton-matrix-zero[simp]}: \text{singleton-matrix } j \ i \ 0 = 0$
 $\langle \text{proof} \rangle$

lemma $\text{nrows-singleton[simp]}: \text{nrows}(\text{singleton-matrix } j \ i \ e) = (\text{if } e = 0 \text{ then } 0 \text{ else } \text{Suc } j)$
 $\langle \text{proof} \rangle$

lemma $\text{ncols-singleton[simp]}: \text{ncols}(\text{singleton-matrix } j \ i \ e) = (\text{if } e = 0 \text{ then } 0 \text{ else } \text{Suc } i)$
 $\langle \text{proof} \rangle$

lemma $\text{combine-singleton}: f \ 0 \ 0 = 0 \implies \text{combine-matrix } f \ (\text{singleton-matrix } j \ i \ a) \ (\text{singleton-matrix } j \ i \ b) = \text{singleton-matrix } j \ i \ (f \ a \ b)$
 $\langle \text{proof} \rangle$

lemma $\text{transpose-singleton[simp]}: \text{transpose-matrix } (\text{singleton-matrix } j \ i \ a) = \text{singleton-matrix } i \ j \ a$
 $\langle \text{proof} \rangle$

lemma $\text{Rep-move-matrix[simp]}:$
 $\text{Rep-matrix } (\text{move-matrix } A \ y \ x) \ j \ i =$
 $(\text{if } (\text{neg } ((\text{int } j) - y)) \mid (\text{neg } ((\text{int } i) - x)) \text{ then } 0 \text{ else } \text{Rep-matrix } A \ (\text{nat}((\text{int } j) - y))$
 $(\text{nat}((\text{int } i) - x)))$
 $\langle \text{proof} \rangle$

lemma $\text{move-matrix-0-0[simp]}: \text{move-matrix } A \ 0 \ 0 = A$
 $\langle \text{proof} \rangle$

lemma $\text{move-matrix-ortho}: \text{move-matrix } A \ j \ i = \text{move-matrix } (\text{move-matrix } A \ j \ 0) \ 0 \ i$

$\langle \text{proof} \rangle$

lemma *transpose-move-matrix*[simp]:

transpose-matrix (*move-matrix* *A* *x* *y*) = *move-matrix* (*transpose-matrix* *A*) *y* *x*
 $\langle \text{proof} \rangle$

lemma *move-matrix-singleton*[simp]: *move-matrix* (*singleton-matrix* *u* *v* *x*) *j* *i* =
(if (*j* + int *u* < 0) | (*i* + int *v* < 0) then 0 else (*singleton-matrix* (nat (*j* + int
u)) (nat (*i* + int *v*)) *x*))
 $\langle \text{proof} \rangle$

lemma *Rep-take-columns*[simp]:

Rep-matrix (*take-columns* *A* *c*) *j* *i* =
(if *i* < *c* then (*Rep-matrix* *A* *j* *i*) else 0)
 $\langle \text{proof} \rangle$

lemma *Rep-take-rows*[simp]:

Rep-matrix (*take-rows* *A* *r*) *j* *i* =
(if *j* < *r* then (*Rep-matrix* *A* *j* *i*) else 0)
 $\langle \text{proof} \rangle$

lemma *Rep-column-of-matrix*[simp]:

Rep-matrix (*column-of-matrix* *A* *c*) *j* *i* = (if *i* = 0 then (*Rep-matrix* *A* *j* *c*) else 0)
 $\langle \text{proof} \rangle$

lemma *Rep-row-of-matrix*[simp]:

Rep-matrix (*row-of-matrix* *A* *r*) *j* *i* = (if *j* = 0 then (*Rep-matrix* *A* *r* *i*) else 0)
 $\langle \text{proof} \rangle$

lemma *column-of-matrix*: *ncols* *A* <= *n* \implies *column-of-matrix* *A* *n* = 0

$\langle \text{proof} \rangle$

lemma *row-of-matrix*: *nrows* *A* <= *n* \implies *row-of-matrix* *A* *n* = 0

$\langle \text{proof} \rangle$

lemma *mult-matrix-singleton-right*[simp]:

assumes *prems*:

! *x*. *fmul* *x* 0 = 0

! *x*. *fmul* 0 *x* = 0

! *x*. *fadd* 0 *x* = *x*

! *x*. *fadd* *x* 0 = *x*

shows (*mult-matrix* *fmul* *fadd* *A* (*singleton-matrix* *j* *i* *e*)) = *apply-matrix* (% *x*.
fmul *x* *e*) (*move-matrix* (*column-of-matrix* *A* *j*) 0 (int *i*))

$\langle \text{proof} \rangle$

lemma *mult-matrix-ext*:

assumes

eprem:

```

    ? e. (! a b. a ≠ b ⟶ fmul a e ≠ fmul b e)
  and fprems:
    ! a. fmul 0 a = 0
    ! a. fmul a 0 = 0
    ! a. fadd a 0 = a
    ! a. fadd 0 a = a
  and contraprems:
    mult-matrix fmul fadd A = mult-matrix fmul fadd B
  shows
    A = B
  ⟨proof⟩

constdefs
  foldmatrix :: ('a ⇒ 'a ⇒ 'a) ⇒ ('a ⇒ 'a ⇒ 'a) ⇒ ('a infmatrix) ⇒ nat ⇒ nat
  ⇒ 'a
  foldmatrix f g A m n == foldseq-transposed g (% j. foldseq f (A j) n) m
  foldmatrix-transposed :: ('a ⇒ 'a ⇒ 'a) ⇒ ('a ⇒ 'a ⇒ 'a) ⇒ ('a infmatrix) ⇒
  nat ⇒ nat ⇒ 'a
  foldmatrix-transposed f g A m n == foldseq g (% j. foldseq-transposed f (A j) n)
  m

lemma foldmatrix-transpose:
  assumes
    ! a b c d. g(f a b) (f c d) = f (g a c) (g b d)
  shows
    foldmatrix f g A m n = foldmatrix-transposed g f (transpose-infmatrix A) n m
  (is ?concl)
  ⟨proof⟩

lemma foldseq-foldseq:
  assumes
    associative f
    associative g
    ! a b c d. g(f a b) (f c d) = f (g a c) (g b d)
  shows
    foldseq g (% j. foldseq f (A j) n) m = foldseq f (% j. foldseq g ((transpose-infmatrix
    A) j) m) n
  ⟨proof⟩

lemma mult-n-nrows:
  assumes
    ! a. fmul 0 a = 0
    ! a. fmul a 0 = 0
    fadd 0 0 = 0
  shows nrows (mult-matrix-n n fmul fadd A B) ≤ nrows A
  ⟨proof⟩

lemma mult-n-ncols:
  assumes

```

! $a. \text{fmul } 0 \ a = 0$
! $a. \text{fmul } a \ 0 = 0$
 $\text{fadd } 0 \ 0 = 0$
shows $\text{ncols } (\text{mult-matrix-n } n \ \text{fmul } \text{fadd } A \ B) \leq \text{ncols } B$
 $\langle \text{proof} \rangle$

lemma *mult-nrows:*

assumes

! $a. \text{fmul } 0 \ a = 0$

! $a. \text{fmul } a \ 0 = 0$

$\text{fadd } 0 \ 0 = 0$

shows $\text{nrows } (\text{mult-matrix } \text{fmul } \text{fadd } A \ B) \leq \text{nrows } A$

$\langle \text{proof} \rangle$

lemma *mult-ncols:*

assumes

! $a. \text{fmul } 0 \ a = 0$

! $a. \text{fmul } a \ 0 = 0$

$\text{fadd } 0 \ 0 = 0$

shows $\text{ncols } (\text{mult-matrix } \text{fmul } \text{fadd } A \ B) \leq \text{ncols } B$

$\langle \text{proof} \rangle$

lemma *nrows-move-matrix-le:* $\text{nrows } (\text{move-matrix } A \ j \ i) \leq \text{nat}((\text{int } (\text{nrows } A)) + j)$

$\langle \text{proof} \rangle$

lemma *ncols-move-matrix-le:* $\text{ncols } (\text{move-matrix } A \ j \ i) \leq \text{nat}((\text{int } (\text{ncols } A)) + i)$

$\langle \text{proof} \rangle$

lemma *mult-matrix-assoc:*

assumes *prems:*

! $a. \text{fmul1 } 0 \ a = 0$

! $a. \text{fmul1 } a \ 0 = 0$

! $a. \text{fmul2 } 0 \ a = 0$

! $a. \text{fmul2 } a \ 0 = 0$

$\text{fadd1 } 0 \ 0 = 0$

$\text{fadd2 } 0 \ 0 = 0$

! $a \ b \ c \ d. \text{fadd2 } (\text{fadd1 } a \ b) (\text{fadd1 } c \ d) = \text{fadd1 } (\text{fadd2 } a \ c) (\text{fadd2 } b \ d)$

associative fadd1

associative fadd2

! $a \ b \ c. \text{fmul2 } (\text{fmul1 } a \ b) \ c = \text{fmul1 } a \ (\text{fmul2 } b \ c)$

! $a \ b \ c. \text{fmul2 } (\text{fadd1 } a \ b) \ c = \text{fadd1 } (\text{fmul2 } a \ c) (\text{fmul2 } b \ c)$

! $a \ b \ c. \text{fmul1 } c \ (\text{fadd2 } a \ b) = \text{fadd2 } (\text{fmul1 } c \ a) (\text{fmul1 } c \ b)$

shows $\text{mult-matrix } \text{fmul2 } \text{fadd2 } (\text{mult-matrix } \text{fmul1 } \text{fadd1 } A \ B) \ C = \text{mult-matrix } \text{fmul1 } \text{fadd1 } A \ (\text{mult-matrix } \text{fmul2 } \text{fadd2 } B \ C)$ (**is** *?concl*)

$\langle \text{proof} \rangle$

lemma

assumes *prems*:
 ! *a*. *fmul1* 0 *a* = 0
 ! *a*. *fmul1* *a* 0 = 0
 ! *a*. *fmul2* 0 *a* = 0
 ! *a*. *fmul2* *a* 0 = 0
fadd1 0 0 = 0
fadd2 0 0 = 0
 ! *a b c d*. *fadd2* (*fadd1* *a b*) (*fadd1* *c d*) = *fadd1* (*fadd2* *a c*) (*fadd2* *b d*)
associative fadd1
associative fadd2
 ! *a b c*. *fmul2* (*fmul1* *a b*) *c* = *fmul1* *a* (*fmul2* *b c*)
 ! *a b c*. *fmul2* (*fadd1* *a b*) *c* = *fadd1* (*fmul2* *a c*) (*fmul2* *b c*)
 ! *a b c*. *fmul1* *c* (*fadd2* *a b*) = *fadd2* (*fmul1* *c a*) (*fmul1* *c b*)
shows
 (*mult-matrix fmul1 fadd1 A*) o (*mult-matrix fmul2 fadd2 B*) = *mult-matrix fmul2 fadd2* (*mult-matrix fmul1 fadd1 A B*)
 <proof>

lemma *mult-matrix-assoc-simple*:
assumes *prems*:
 ! *a*. *fmul* 0 *a* = 0
 ! *a*. *fmul* *a* 0 = 0
fadd 0 0 = 0
associative fadd
commutative fadd
associative fmul
distributive fmul fadd
shows *mult-matrix fmul fadd* (*mult-matrix fmul fadd A B*) *C* = *mult-matrix fmul fadd A* (*mult-matrix fmul fadd B C*) (**is** ?concl)
 <proof>

lemma *transpose-apply-matrix*: *f* 0 = 0 \implies *transpose-matrix* (*apply-matrix f A*)
 = *apply-matrix f* (*transpose-matrix A*)
 <proof>

lemma *transpose-combine-matrix*: *f* 0 0 = 0 \implies *transpose-matrix* (*combine-matrix f A B*)
 = *combine-matrix f* (*transpose-matrix A*) (*transpose-matrix B*)
 <proof>

lemma *Rep-mult-matrix*:
assumes
 ! *a*. *fmul* 0 *a* = 0
 ! *a*. *fmul* *a* 0 = 0
fadd 0 0 = 0
shows
Rep-matrix(*mult-matrix fmul fadd A B*) *j i* =
 foldseq *fadd* (% *k*. *fmul* (*Rep-matrix A j k*) (*Rep-matrix B k i*)) (*max* (*ncols A*)
 (*nrows B*))
 <proof>

lemma *transpose-mult-matrix*:

assumes

$! a. \text{fmul } 0 \ a = 0$

$! a. \text{fmul } a \ 0 = 0$

$\text{fadd } 0 \ 0 = 0$

$! x \ y. \text{fmul } y \ x = \text{fmul } x \ y$

shows

$\text{transpose-matrix } (\text{mult-matrix } \text{fmul } \text{fadd } A \ B) = \text{mult-matrix } \text{fmul } \text{fadd } (\text{transpose-matrix } B) \ (\text{transpose-matrix } A)$
 $\langle \text{proof} \rangle$

lemma *column-transpose-matrix*: $\text{column-of-matrix } (\text{transpose-matrix } A) \ n = \text{transpose-matrix } (\text{row-of-matrix } A \ n)$
 $\langle \text{proof} \rangle$

lemma *take-columns-transpose-matrix*: $\text{take-columns } (\text{transpose-matrix } A) \ n = \text{transpose-matrix } (\text{take-rows } A \ n)$
 $\langle \text{proof} \rangle$

instance *matrix* :: $(\{\text{ord}, \text{zero}\}) \ \text{ord} \ \langle \text{proof} \rangle$

defs (**overloaded**)

$\text{le-matrix-def}: (A :: ('a :: \{\text{ord}, \text{zero}\}) \ \text{matrix}) \leq B == ! j \ i. (\text{Rep-matrix } A \ j \ i) \leq (\text{Rep-matrix } B \ j \ i)$

$\text{less-def}: (A :: ('a :: \{\text{ord}, \text{zero}\}) \ \text{matrix}) < B == (A \leq B) \ \& \ (A \neq B)$

instance *matrix* :: $(\{\text{order}, \text{zero}\}) \ \text{order} \ \langle \text{proof} \rangle$

lemma *le-apply-matrix*:

assumes

$f \ 0 = 0$

$! x \ y. x \leq y \longrightarrow f \ x \leq f \ y$

$(a :: ('a :: \{\text{ord}, \text{zero}\}) \ \text{matrix}) \leq b$

shows

$\text{apply-matrix } f \ a \leq \text{apply-matrix } f \ b$

$\langle \text{proof} \rangle$

lemma *le-combine-matrix*:

assumes

$f \ 0 \ 0 = 0$

$! a \ b \ c \ d. a \leq b \ \& \ c \leq d \longrightarrow f \ a \ c \leq f \ b \ d$

$A \leq B$

$C \leq D$

shows

$\text{combine-matrix } f \ A \ C \leq \text{combine-matrix } f \ B \ D$

$\langle \text{proof} \rangle$

lemma *le-left-combine-matrix*:

assumes

$f\ 0\ 0 = 0$

$! a\ b\ c. a \leq b \longrightarrow f\ c\ a \leq f\ c\ b$

$A \leq B$

shows

$combine_matrix\ f\ C\ A \leq combine_matrix\ f\ C\ B$

$\langle proof \rangle$

lemma *le-right-combine-matrix*:

assumes

$f\ 0\ 0 = 0$

$! a\ b\ c. a \leq b \longrightarrow f\ a\ c \leq f\ b\ c$

$A \leq B$

shows

$combine_matrix\ f\ A\ C \leq combine_matrix\ f\ B\ C$

$\langle proof \rangle$

lemma *le-transpose-matrix*: $(A \leq B) = (transpose_matrix\ A \leq transpose_matrix\ B)$

$\langle proof \rangle$

lemma *le-foldseq*:

assumes

$! a\ b\ c\ d. a \leq b \ \& \ c \leq d \longrightarrow f\ a\ c \leq f\ b\ d$

$! i. i \leq n \longrightarrow s\ i \leq t\ i$

shows

$foldseq\ f\ s\ n \leq foldseq\ f\ t\ n$

$\langle proof \rangle$

lemma *le-left-mult*:

assumes

$! a\ b\ c\ d. a \leq b \ \& \ c \leq d \longrightarrow fadd\ a\ c \leq fadd\ b\ d$

$! c\ a\ b. 0 \leq c \ \& \ a \leq b \longrightarrow fmul\ c\ a \leq fmul\ c\ b$

$! a. fmul\ 0\ a = 0$

$! a. fmul\ a\ 0 = 0$

$fadd\ 0\ 0 = 0$

$0 \leq C$

$A \leq B$

shows

$mult_matrix\ fmul\ fadd\ C\ A \leq mult_matrix\ fmul\ fadd\ C\ B$

$\langle proof \rangle$

lemma *le-right-mult*:

assumes

$! a\ b\ c\ d. a \leq b \ \& \ c \leq d \longrightarrow fadd\ a\ c \leq fadd\ b\ d$

$! c\ a\ b. 0 \leq c \ \& \ a \leq b \longrightarrow fmul\ a\ c \leq fmul\ b\ c$

$! a. fmul\ 0\ a = 0$

$! a. fmul\ a\ 0 = 0$


```

fadd 0 0 = 0
0 <= C
A <= B
shows
mult-matrix fmul fadd A C <= mult-matrix fmul fadd B C
⟨proof⟩

lemma spec2: ! j i. P j i ⟹ P j i ⟨proof⟩
lemma neg-imp: (¬ Q ⟹ ¬ P) ⟹ P ⟹ Q ⟨proof⟩

lemma singleton-matrix-le[simp]: (singleton-matrix j i a <= singleton-matrix j i
b) = (a <= (b::order))
⟨proof⟩

lemma singleton-le-zero[simp]: (singleton-matrix j i x <= 0) = (x <= (0::'a::{order,zero}))
⟨proof⟩

lemma singleton-ge-zero[simp]: (0 <= singleton-matrix j i x) = ((0::'a::{order,zero})
<= x)
⟨proof⟩

lemma move-matrix-le-zero[simp]: 0 <= j ⟹ 0 <= i ⟹ (move-matrix A j i
<= 0) = (A <= (0::('a::{order,zero}) matrix))
⟨proof⟩

lemma move-matrix-zero-le[simp]: 0 <= j ⟹ 0 <= i ⟹ (0 <= move-matrix
A j i) = ((0::('a::{order,zero}) matrix) <= A)
⟨proof⟩

lemma move-matrix-le-move-matrix-iff[simp]: 0 <= j ⟹ 0 <= i ⟹ (move-matrix
A j i <= move-matrix B j i) = (A <= (B::('a::{order,zero}) matrix))
⟨proof⟩

end

theory Matrix=MatrixGeneral:

instance matrix :: (minus) minus
⟨proof⟩

instance matrix :: (plus) plus
⟨proof⟩

instance matrix :: ({plus,times}) times
⟨proof⟩

defs (overloaded)

```

```

plus-matrix-def:  $A + B == \text{combine-matrix } (op +) A B$ 
diff-matrix-def:  $A - B == \text{combine-matrix } (op -) A B$ 
minus-matrix-def:  $- A == \text{apply-matrix } \text{uminus } A$ 
times-matrix-def:  $A * B == \text{mult-matrix } (op *) (op +) A B$ 

lemma is-meet-combine-matrix-meet: is-meet (combine-matrix meet)
  ⟨proof⟩

lemma is-join-combine-matrix-join: is-join (combine-matrix join)
  ⟨proof⟩

instance matrix :: (ordered-ab-group) ordered-ab-group-meet
  ⟨proof⟩

defs (overloaded)
  abs-matrix-def: abs ( $A :: ('a :: \text{ordered-ab-group}) \text{matrix}$ ) == join  $A$  ( $- A$ )

instance matrix :: (ordered-ring) ordered-ring
  ⟨proof⟩

lemma Rep-matrix-add[simp]: Rep-matrix (( $a :: ('a :: \text{ordered-ab-group}) \text{matrix}$ ) +  $b$ )
   $j \ i = (\text{Rep-matrix } a \ j \ i) + (\text{Rep-matrix } b \ j \ i)$ 
  ⟨proof⟩

lemma Rep-matrix-mult: Rep-matrix (( $a :: ('a :: \text{ordered-ring}) \text{matrix}$ ) *  $b$ )  $j \ i =$ 
  foldseq (op +) (%  $k.$  (Rep-matrix  $a \ j \ k$ ) * (Rep-matrix  $b \ k \ i$ )) (max (ncols  $a$ )
  (nrows  $b$ ))
  ⟨proof⟩

lemma apply-matrix-add: !  $x \ y. f (x+y) = (f x) + (f y) \implies f 0 = (0 :: 'a) \implies$ 
  apply-matrix  $f$  (( $a :: ('a :: \text{ordered-ab-group}) \text{matrix}$ ) +  $b$ ) = (apply-matrix  $f$   $a$ ) +
  (apply-matrix  $f$   $b$ )
  ⟨proof⟩

lemma singleton-matrix-add: singleton-matrix  $j \ i$  (( $a :: (\text{ordered-ab-group})$ ) +  $b$ ) =
  (singleton-matrix  $j \ i$   $a$ ) + (singleton-matrix  $j \ i$   $b$ )
  ⟨proof⟩

lemma nrows-mult: nrows (( $A :: ('a :: \text{ordered-ring}) \text{matrix}$ ) *  $B$ ) <= nrows  $A$ 
  ⟨proof⟩

lemma ncols-mult: ncols (( $A :: ('a :: \text{ordered-ring}) \text{matrix}$ ) *  $B$ ) <= ncols  $B$ 
  ⟨proof⟩

constdefs
  one-matrix :: nat  $\Rightarrow ('a :: \{zero, one\}) \text{matrix}$ 
  one-matrix  $n == \text{Abs-matrix } (\% \ j \ i. \text{if } j = i \ \& \ j < n \text{ then } 1 \text{ else } 0)$ 

```

lemma *Rep-one-matrix[simp]*: *Rep-matrix* (*one-matrix* *n*) *j i* = (if (*j* = *i* & *j* < *n*) then 1 else 0)

⟨proof⟩

lemma *nrows-one-matrix[simp]*: *nrows* ((*one-matrix* *n*) :: ('a::axclass-0-neq-1)matrix) = *n* (is ?r = -)

⟨proof⟩

lemma *ncols-one-matrix[simp]*: *ncols* ((*one-matrix* *n*) :: ('a::axclass-0-neq-1)matrix) = *n* (is ?r = -)

⟨proof⟩

lemma *one-matrix-mult-right[simp]*: *ncols* *A* ≤ *n* ⇒ (*A*::('a::{lordered-ring, ring-1})matrix) * (*one-matrix* *n*) = *A*

⟨proof⟩

lemma *one-matrix-mult-left[simp]*: *nrows* *A* ≤ *n* ⇒ (*one-matrix* *n*) * *A* = (*A*::('a::{lordered-ring, ring-1})matrix)

⟨proof⟩

lemma *transpose-matrix-mult*: *transpose-matrix* ((*A*::('a::{lordered-ring, comm-ring})matrix)**B*) = (*transpose-matrix* *B*) * (*transpose-matrix* *A*)

⟨proof⟩

lemma *transpose-matrix-add*: *transpose-matrix* ((*A*::('a::lordered-ab-group)matrix)+*B*) = *transpose-matrix* *A* + *transpose-matrix* *B*

⟨proof⟩

lemma *transpose-matrix-diff*: *transpose-matrix* ((*A*::('a::lordered-ab-group)matrix)−*B*) = *transpose-matrix* *A* − *transpose-matrix* *B*

⟨proof⟩

lemma *transpose-matrix-minus*: *transpose-matrix* (−(*A*::('a::lordered-ring)matrix)) = − *transpose-matrix* (*A*::('a::lordered-ring)matrix)

⟨proof⟩

constdefs

right-inverse-matrix :: ('a::{lordered-ring, ring-1})matrix ⇒ 'a matrix ⇒ bool
right-inverse-matrix *A X* == (*A* * *X* = *one-matrix* (max (*nrows* *A*) (*ncols* *X*)))

∧ *nrows* *X* ≤ *ncols* *A*

left-inverse-matrix :: ('a::{lordered-ring, ring-1})matrix ⇒ 'a matrix ⇒ bool
left-inverse-matrix *A X* == (*X* * *A* = *one-matrix* (max(*nrows* *X*) (*ncols* *A*))) ∧

ncols *X* ≤ *nrows* *A*

inverse-matrix :: ('a::{lordered-ring, ring-1})matrix ⇒ 'a matrix ⇒ bool

inverse-matrix *A X* == (*right-inverse-matrix* *A X*) ∧ (*left-inverse-matrix* *A X*)

lemma *right-inverse-matrix-dim*: *right-inverse-matrix* *A X* ⇒ *nrows* *A* = *ncols* *X*

⟨proof⟩

lemma *left-inverse-matrix-dim*: *left-inverse-matrix* A $Y \implies \text{ncols } A = \text{nrows } Y$
 <proof>

lemma *left-right-inverse-matrix-unique*:
 assumes *left-inverse-matrix* A Y *right-inverse-matrix* A X
 shows $X = Y$
 <proof>

lemma *inverse-matrix-inject*: $\llbracket \text{inverse-matrix } A \ X; \text{inverse-matrix } A \ Y \rrbracket \implies X = Y$
 <proof>

lemma *one-matrix-inverse*: *inverse-matrix* (*one-matrix* n) (*one-matrix* n)
 <proof>

lemma *zero-imp-mult-zero*: $(a::'a::\text{ring}) = 0 \mid b = 0 \implies a * b = 0$
 <proof>

lemma *Rep-matrix-zero-imp-mult-zero*:
 ! $j \ i \ k. (\text{Rep-matrix } A \ j \ k = 0) \mid (\text{Rep-matrix } B \ k \ i) = 0 \implies A * B =$
 $(0::('a::\text{lordered-ring}) \text{ matrix})$
 <proof>

lemma *add-nrows*: $\text{nrows } (A::('a::\text{comm-monoid-add}) \text{ matrix}) \leq u \implies \text{nrows } B \leq u \implies \text{nrows } (A + B) \leq u$
 <proof>

lemma *move-matrix-row-mult*: *move-matrix* $((A::('a::\text{lordered-ring}) \text{ matrix}) * B)$
 $j \ 0 = (\text{move-matrix } A \ j \ 0) * B$
 <proof>

lemma *move-matrix-col-mult*: *move-matrix* $((A::('a::\text{lordered-ring}) \text{ matrix}) * B)$
 $0 \ i = A * (\text{move-matrix } B \ 0 \ i)$
 <proof>

lemma *move-matrix-add*: $((\text{move-matrix } (A + B) \ j \ i)::('a::\text{lordered-ab-group}) \text{ matrix})) = (\text{move-matrix } A \ j \ i) + (\text{move-matrix } B \ j \ i)$
 <proof>

lemma *move-matrix-mult*: *move-matrix* $((A::('a::\text{lordered-ring}) \text{ matrix}) * B) \ j \ i =$
 $(\text{move-matrix } A \ j \ 0) * (\text{move-matrix } B \ 0 \ i)$
 <proof>

constdefs
scalar-mult :: $('a::\text{lordered-ring}) \Rightarrow 'a \text{ matrix} \Rightarrow 'a \text{ matrix}$
scalar-mult $a \ m == \text{apply-matrix } (\text{op } * \ a) \ m$

lemma *scalar-mult-zero[simp]*: *scalar-mult* $y \ 0 = 0$

```

    <proof>

lemma scalar-mult-add: scalar-mult y (a+b) = (scalar-mult y a) + (scalar-mult y
b)
    <proof>

lemma Rep-scalar-mult[simp]: Rep-matrix (scalar-mult y a) j i = y * (Rep-matrix
a j i)
    <proof>

lemma scalar-mult-singleton[simp]: scalar-mult y (singleton-matrix j i x) = singleton-matrix
j i (y * x)
    <proof>

lemma Rep-minus[simp]: Rep-matrix (-(A:::ordered-ab-group)) x y = - (Rep-matrix
A x y)
    <proof>

lemma join-matrix: join (A::('a::ordered-ring) matrix) B = combine-matrix join
A B
    <proof>

lemma meet-matrix: meet (A::('a::ordered-ring) matrix) B = combine-matrix
meet A B
    <proof>

lemma Rep-abs[simp]: Rep-matrix (abs (A:::ordered-ring)) x y = abs (Rep-matrix
A x y)
    <proof>

end

theory SparseMatrix imports Matrix begin

types
    'a svec = (nat * 'a) list
    'a smat = ('a svec) svec

consts
    sparse-row-vector :: ('a::ordered-ring) svec  $\Rightarrow$  'a matrix
    sparse-row-matrix :: ('a::ordered-ring) smat  $\Rightarrow$  'a matrix

defs
    sparse-row-vector-def : sparse-row-vector arr == foldl (% m x. m + (singleton-matrix
0 (fst x) (snd x))) 0 arr
    sparse-row-matrix-def : sparse-row-matrix arr == foldl (% m r. m + (move-matrix
(sparse-row-vector (snd r)) (int (fst r)) 0)) 0 arr

```

lemma *sparse-row-vector-empty*[simp]: *sparse-row-vector* [] = 0
 ⟨proof⟩

lemma *sparse-row-matrix-empty*[simp]: *sparse-row-matrix* [] = 0
 ⟨proof⟩

lemma *foldl-distrstart*[rule-format]: ! *a x y*. (*f* (*g x y*) *a* = *g x* (*f y a*)) \implies ! *x y*.
 (*foldl f* (*g x y*) *l* = *g x* (*foldl f y l*))
 ⟨proof⟩

lemma *sparse-row-vector-cons*[simp]: *sparse-row-vector* (*a* # *arr*) = (*singleton-matrix* 0 (*fst a*) (*snd a*)) + (*sparse-row-vector* *arr*)
 ⟨proof⟩

lemma *sparse-row-vector-append*[simp]: *sparse-row-vector* (*a* @ *b*) = (*sparse-row-vector* *a*) + (*sparse-row-vector* *b*)
 ⟨proof⟩

lemma *nrows-spvec*[simp]: *nrows* (*sparse-row-vector* *x*) <= (*Suc* 0)
 ⟨proof⟩

lemma *sparse-row-matrix-cons*: *sparse-row-matrix* (*a* # *arr*) = ((*move-matrix* (*sparse-row-vector* (*snd a*)) (*int* (*fst a*)) 0)) + *sparse-row-matrix* *arr*
 ⟨proof⟩

lemma *sparse-row-matrix-append*: *sparse-row-matrix* (*arr* @ *brr*) = (*sparse-row-matrix* *arr*) + (*sparse-row-matrix* *brr*)
 ⟨proof⟩

consts

sorted-spvec :: 'a *spvec* \Rightarrow bool
sorted-spmat :: 'a *spmat* \Rightarrow bool

primrec

sorted-spmat [] = True
sorted-spmat (*a* # *as*) = ((*sorted-spvec* (*snd a*)) & (*sorted-spmat* *as*))

primrec

sorted-spvec [] = True
sorted-spvec-step: *sorted-spvec* (*a* # *as*) = (*case as of* [] \Rightarrow True | *b* # *bs* \Rightarrow ((*fst a* < *fst b*) & (*sorted-spvec* *as*)))

declare *sorted-spvec.simps* [simp del]

lemma *sorted-spvec-empty*[simp]: *sorted-spvec* [] = True
 ⟨proof⟩

lemma *sorted-spvec-cons1*: *sorted-spvec* (*a* # *as*) \implies *sorted-spvec* *as*

$\langle \text{proof} \rangle$

lemma *sorted-spvec-cons2*: $\text{sorted-spvec } (a \# b \# t) \implies \text{sorted-spvec } (a \# t)$
 $\langle \text{proof} \rangle$

lemma *sorted-spvec-cons3*: $\text{sorted-spvec}(a \# b \# t) \implies \text{fst } a < \text{fst } b$
 $\langle \text{proof} \rangle$

lemma *sorted-sparse-row-vector-zero*[*rule-format*]: $m \leq n \longrightarrow \text{sorted-spvec } ((n, a) \# \text{arr})$
 $\longrightarrow \text{Rep-matrix } (\text{sparse-row-vector } \text{arr}) \ j \ m = 0$
 $\langle \text{proof} \rangle$

lemma *sorted-sparse-row-matrix-zero*[*rule-format*]: $m \leq n \longrightarrow \text{sorted-spvec } ((n, a) \# \text{arr})$
 $\longrightarrow \text{Rep-matrix } (\text{sparse-row-matrix } \text{arr}) \ m \ j = 0$
 $\langle \text{proof} \rangle$

consts

$\text{abs-spvec} :: ('a :: \text{lordered-ring}) \text{ spvec} \Rightarrow 'a \text{ spvec}$
 $\text{minus-spvec} :: ('a :: \text{lordered-ring}) \text{ spvec} \Rightarrow 'a \text{ spvec}$
 $\text{smult-spvec} :: ('a :: \text{lordered-ring}) \Rightarrow 'a \text{ spvec} \Rightarrow 'a \text{ spvec}$
 $\text{addmult-spvec} :: ('a :: \text{lordered-ring}) * 'a \text{ spvec} * 'a \text{ spvec} \Rightarrow 'a \text{ spvec}$

primrec

$\text{minus-spvec } [] = []$
 $\text{minus-spvec } (a \# as) = (\text{fst } a, -(\text{snd } a)) \# (\text{minus-spvec } as)$

primrec

$\text{abs-spvec } [] = []$
 $\text{abs-spvec } (a \# as) = (\text{fst } a, \text{abs } (\text{snd } a)) \# (\text{abs-spvec } as)$

lemma *sparse-row-vector-minus*:

$\text{sparse-row-vector } (\text{minus-spvec } v) = - (\text{sparse-row-vector } v)$
 $\langle \text{proof} \rangle$

lemma *sparse-row-vector-abs*:

$\text{sorted-spvec } v \implies \text{sparse-row-vector } (\text{abs-spvec } v) = \text{abs } (\text{sparse-row-vector } v)$
 $\langle \text{proof} \rangle$

lemma *sorted-spvec-minus-spvec*:

$\text{sorted-spvec } v \implies \text{sorted-spvec } (\text{minus-spvec } v)$
 $\langle \text{proof} \rangle$

lemma *sorted-spvec-minus-spvec*:

$\text{sorted-spvec } v \implies \text{sorted-spvec } (\text{minus-spvec } v)$
 $\langle \text{proof} \rangle$

lemma *sorted-spvec-abs-spvec*:

$\text{sorted-spvec } v \implies \text{sorted-spvec } (\text{abs-spvec } v)$
 $\langle \text{proof} \rangle$

defs

smult-spvec-def: $smult\text{-}spvec\ y\ arr == map\ (\% a.\ (fst\ a,\ y * snd\ a))\ arr$

lemma *smult-spvec-empty[simp]*: $smult\text{-}spvec\ y\ [] = []$
 ⟨proof⟩

lemma *smult-spvec-cons*: $smult\text{-}spvec\ y\ (a\#arr) = (fst\ a,\ y * (snd\ a))\ \# (smult\text{-}spvec\ y\ arr)$
 ⟨proof⟩

recdef *addmult-spvec measure* (% (y, a, b). length a + (length b))
 addmult-spvec (y, arr, []) = arr
 addmult-spvec (y, [], brr) = smult-spvec y brr
 addmult-spvec (y, a#arr, b#brr) = (
 if (fst a) < (fst b) then (a#(addmult-spvec (y, arr, b#brr)))
 else (if (fst b < fst a) then ((fst b, y * (snd b))#(addmult-spvec (y, a#arr, brr)))
 else ((fst a, (snd a) + y*(snd b))#(addmult-spvec (y, arr, brr))))

lemma *addmult-spvec-empty1[simp]*: $addmult\text{-}spvec\ (y,\ [],\ a) = smult\text{-}spvec\ y\ a$
 ⟨proof⟩

lemma *addmult-spvec-empty2[simp]*: $addmult\text{-}spvec\ (y,\ a,\ []) = a$
 ⟨proof⟩

lemma *sparse-row-vector-map*: $(! x\ y.\ f\ (x+y) = (f\ x) + (f\ y)) \implies (f::'a \Rightarrow ('a::lordered\text{-}ring))$
 $0 = 0 \implies$
 sparse-row-vector (map (% x. (fst x, f (snd x))) a) = apply-matrix f (sparse-row-vector a)
 ⟨proof⟩

lemma *sparse-row-vector-smult*: $sparse\text{-}row\text{-}vector\ (smult\text{-}spvec\ y\ a) = scalar\text{-}mult\ y\ (sparse\text{-}row\text{-}vector\ a)$
 ⟨proof⟩

lemma *sparse-row-vector-addmult-spvec*: $sparse\text{-}row\text{-}vector\ (addmult\text{-}spvec\ (y::'a::lordered\text{-}ring,\ a,\ b)) =$
 $(sparse\text{-}row\text{-}vector\ a) + (scalar\text{-}mult\ y\ (sparse\text{-}row\text{-}vector\ b))$
 ⟨proof⟩

lemma *sorted-smult-spvec[rule-format]*: $sorted\text{-}spvec\ a \implies sorted\text{-}spvec\ (smult\text{-}spvec\ y\ a)$
 ⟨proof⟩

lemma *sorted-spvec-addmult-spvec-helper*: $\llbracket sorted\text{-}spvec\ (addmult\text{-}spvec\ (y,\ (a,\ b)\ \# arr,\ brr));\ aa < a;\ sorted\text{-}spvec\ ((a,\ b)\ \# arr);\ sorted\text{-}spvec\ ((aa,\ ba)\ \# brr) \rrbracket \implies sorted\text{-}spvec\ ((aa,\ y * ba)\ \# addmult\text{-}spvec\ (y,\ (a,\ b)\ \# arr,\ brr))$

$\langle \text{proof} \rangle$

lemma *sorted-spvec-addmult-spvec-helper2*:

$\llbracket \text{sorted-spvec } (\text{addmult-spvec } (y, \text{arr}, (aa, ba) \# \text{brr})); a < aa; \text{sorted-spvec } ((a, b) \# \text{arr}); \text{sorted-spvec } ((aa, ba) \# \text{brr}) \rrbracket$
 $\implies \text{sorted-spvec } ((a, b) \# \text{addmult-spvec } (y, \text{arr}, (aa, ba) \# \text{brr}))$
 $\langle \text{proof} \rangle$

lemma *sorted-spvec-addmult-spvec-helper3*[*rule-format*]:

$\text{sorted-spvec } (\text{addmult-spvec } (y, \text{arr}, \text{brr})) \longrightarrow \text{sorted-spvec } ((aa, b) \# \text{arr}) \longrightarrow$
 $\text{sorted-spvec } ((aa, ba) \# \text{brr})$
 $\longrightarrow \text{sorted-spvec } ((aa, b + y * ba) \# (\text{addmult-spvec } (y, \text{arr}, \text{brr})))$
 $\langle \text{proof} \rangle$

lemma *sorted-addmult-spvec*[*rule-format*]: $\text{sorted-spvec } a \longrightarrow \text{sorted-spvec } b \longrightarrow$
 $\text{sorted-spvec } (\text{addmult-spvec } (y, a, b))$
 $\langle \text{proof} \rangle$

consts

$\text{mult-spvec-spmat} :: ('a::\text{lordered-ring}) \text{ spvec} * 'a \text{ spvec} * 'a \text{ smat} \Rightarrow 'a \text{ spvec}$

recdef *mult-spvec-spmat measure* (% (c, arr, brr). (length arr) + (length brr))

$\text{mult-spvec-spmat } (c, [], \text{brr}) = c$
 $\text{mult-spvec-spmat } (c, \text{arr}, []) = c$
 $\text{mult-spvec-spmat } (c, a \# \text{arr}, b \# \text{brr}) =$
 $\text{if } ((\text{fst } a) < (\text{fst } b)) \text{ then } (\text{mult-spvec-spmat } (c, \text{arr}, b \# \text{brr}))$
 $\text{else if } ((\text{fst } b) < (\text{fst } a)) \text{ then } (\text{mult-spvec-spmat } (c, a \# \text{arr}, \text{brr}))$
 $\text{else } (\text{mult-spvec-spmat } (\text{addmult-spvec } (\text{snd } a, c, \text{snd } b), \text{arr}, \text{brr})))$

lemma *sparse-row-mult-spvec-spmat*[*rule-format*]: $\text{sorted-spvec } (a::('a::\text{lordered-ring})$
 $\text{spvec}) \longrightarrow \text{sorted-spvec } B \longrightarrow$

$\text{sparse-row-vector } (\text{mult-spvec-spmat } (c, a, B)) = (\text{sparse-row-vector } c) + (\text{sparse-row-vector } a) * (\text{sparse-row-matrix } B)$
 $\langle \text{proof} \rangle$

lemma *sorted-mult-spvec-spmat*[*rule-format*]:

$\text{sorted-spvec } (c::('a::\text{lordered-ring}) \text{ spvec}) \longrightarrow \text{sorted-spmat } B \longrightarrow \text{sorted-spvec}$
 $(\text{mult-spvec-spmat } (c, a, B))$
 $\langle \text{proof} \rangle$

consts

$\text{mult-spmat} :: ('a::\text{lordered-ring}) \text{ smat} \Rightarrow 'a \text{ smat} \Rightarrow 'a \text{ smat}$

primrec

$\text{mult-spmat } [] \ A = []$
 $\text{mult-spmat } (a \# as) \ A = (\text{fst } a, \text{mult-spvec-spmat } ([], \text{snd } a, A)) \# (\text{mult-spmat } as \ A)$

lemma *sparse-row-mult-spmat*[*rule-format*]:

$sorted\text{-}spmat\ A \longrightarrow sorted\text{-}spvec\ B \longrightarrow sparse\text{-}row\text{-}matrix\ (mult\text{-}spmat\ A\ B) =$
 $(sparse\text{-}row\text{-}matrix\ A) * (sparse\text{-}row\text{-}matrix\ B)$
 $\langle proof \rangle$

lemma *sorted-spvec-mult-spmat*[rule-format]:
 $sorted\text{-}spvec\ (A::('a::ordered\text{-}ring)\ spmat) \longrightarrow sorted\text{-}spvec\ (mult\text{-}spmat\ A\ B)$
 $\langle proof \rangle$

lemma *sorted-spmat-mult-spmat*[rule-format]:
 $sorted\text{-}spmat\ (B::('a::ordered\text{-}ring)\ spmat) \longrightarrow sorted\text{-}spmat\ (mult\text{-}spmat\ A\ B)$
 $\langle proof \rangle$

consts

$add\text{-}spvec :: ('a::ordered\text{-}ab\text{-}group)\ spvec * 'a\ spvec \Rightarrow 'a\ spvec$
 $add\text{-}spmat :: ('a::ordered\text{-}ab\text{-}group)\ spmat * 'a\ spmat \Rightarrow 'a\ spmat$

recdef *add-spvec measure* (% (a, b). length a + (length b))
 $add\text{-}spvec\ (arr, []) = arr$
 $add\text{-}spvec\ ([], brr) = brr$
 $add\text{-}spvec\ (a\#arr, b\#brr) =$
 $if\ (fst\ a) < (fst\ b)\ then\ (a\#(add\text{-}spvec\ (arr, b\#brr)))$
 $else\ (if\ (fst\ b < fst\ a)\ then\ (b\#(add\text{-}spvec\ (a\#arr, brr)))$
 $else\ ((fst\ a, (snd\ a)+(snd\ b))\#(add\text{-}spvec\ (arr, brr))))$

lemma *add-spvec-empty1*[simp]: $add\text{-}spvec\ ([], a) = a$
 $\langle proof \rangle$

lemma *add-spvec-empty2*[simp]: $add\text{-}spvec\ (a, []) = a$
 $\langle proof \rangle$

lemma *sparse-row-vector-add*: $sparse\text{-}row\text{-}vector\ (add\text{-}spvec\ (a, b)) = (sparse\text{-}row\text{-}vector\ a) + (sparse\text{-}row\text{-}vector\ b)$
 $\langle proof \rangle$

recdef *add-spmat measure* (% (A,B). (length A)+(length B))
 $add\text{-}spmat\ ([], bs) = bs$
 $add\text{-}spmat\ (as, []) = as$
 $add\text{-}spmat\ (a\#as, b\#bs) =$
 $if\ fst\ a < fst\ b\ then$
 $(a\#(add\text{-}spmat\ (as, b\#bs)))$
 $else\ (if\ fst\ b < fst\ a\ then$
 $(b\#(add\text{-}spmat\ (a\#as, bs)))$
 $else$
 $((fst\ a, add\text{-}spvec\ (snd\ a, snd\ b))\#(add\text{-}spmat\ (as, bs))))$

lemma *sparse-row-add-spmat*: $sparse\text{-}row\text{-}matrix\ (add\text{-}spmat\ (A, B)) = (sparse\text{-}row\text{-}matrix\ A) + (sparse\text{-}row\text{-}matrix\ B)$
 $\langle proof \rangle$

lemma *sorted-add-spvec-helper1*[rule-format]: $\text{add-spvec } ((a,b)\#arr, brr) = (ab, bb) \# list \longrightarrow (ab = a \mid (brr \neq [] \ \& \ ab = fst \ (hd \ brr)))$
 <proof>

lemma *sorted-add-spmat-helper1*[rule-format]: $\text{add-spmat } ((a,b)\#arr, brr) = (ab, bb) \# list \longrightarrow (ab = a \mid (brr \neq [] \ \& \ ab = fst \ (hd \ brr)))$
 <proof>

lemma *sorted-add-spvec-helper*[rule-format]: $\text{add-spvec } (arr, brr) = (ab, bb) \# list \longrightarrow ((arr \neq [] \ \& \ ab = fst \ (hd \ arr)) \mid (brr \neq [] \ \& \ ab = fst \ (hd \ brr)))$
 <proof>

lemma *sorted-add-spmat-helper*[rule-format]: $\text{add-spmat } (arr, brr) = (ab, bb) \# list \longrightarrow ((arr \neq [] \ \& \ ab = fst \ (hd \ arr)) \mid (brr \neq [] \ \& \ ab = fst \ (hd \ brr)))$
 <proof>

lemma *add-spvec-commute*: $\text{add-spvec } (a, b) = \text{add-spvec } (b, a)$
 <proof>

lemma *add-spmat-commute*: $\text{add-spmat } (a, b) = \text{add-spmat } (b, a)$
 <proof>

lemma *sorted-add-spvec-helper2*: $\text{add-spvec } ((a,b)\#arr, brr) = (ab, bb) \# list \implies aa < a \implies \text{sorted-spvec } ((aa, ba) \# brr) \implies aa < ab$
 <proof>

lemma *sorted-add-spmat-helper2*: $\text{add-spmat } ((a,b)\#arr, brr) = (ab, bb) \# list \implies aa < a \implies \text{sorted-spmat } ((aa, ba) \# brr) \implies aa < ab$
 <proof>

lemma *sorted-spvec-add-spvec*[rule-format]: $\text{sorted-spvec } a \longrightarrow \text{sorted-spvec } b \longrightarrow \text{sorted-spvec } (\text{add-spvec } (a, b))$
 <proof>

lemma *sorted-spvec-add-spmat*[rule-format]: $\text{sorted-spvec } A \longrightarrow \text{sorted-spvec } B \longrightarrow \text{sorted-spmat } (\text{add-spmat } (A, B))$
 <proof>

lemma *sorted-spmat-add-spmat*[rule-format]: $\text{sorted-spmat } A \longrightarrow \text{sorted-spmat } B \longrightarrow \text{sorted-spmat } (\text{add-spmat } (A, B))$
 <proof>

consts

$\text{le-spvec} :: ('a::\text{ordered-ab-group}) \text{ spvec} * 'a \text{ spvec} \Rightarrow \text{bool}$
 $\text{le-spmat} :: ('a::\text{ordered-ab-group}) \text{ spat} * 'a \text{ spat} \Rightarrow \text{bool}$

recdef *le-spvec measure* (% (a,b). (length a) + (length b))
 $\text{le-spvec } ([], []) = \text{True}$
 $\text{le-spvec } (a\#as, []) = ((snd \ a \leq 0) \ \& \ (\text{le-spvec } (as, [])))$

```

le-spvec ([], b#bs) = ((0 <= snd b) & (le-spvec ([], bs)))
le-spvec (a#as, b#bs) = (
  if (fst a < fst b) then
    ((snd a <= 0) & (le-spvec (as, b#bs)))
  else (if (fst b < fst a) then
    ((0 <= snd b) & (le-spvec (a#as, bs)))
  else
    ((snd a <= snd b) & (le-spvec (as, bs))))))

recdef le-spmat measure (% (a,b). (length a) + (length b))
le-spmat ([], []) = True
le-spmat (a#as, []) = (le-spvec (snd a, []) & (le-spmat (as, [])))
le-spmat ([], b#bs) = (le-spvec ([], snd b) & (le-spmat ([], bs)))
le-spmat (a#as, b#bs) = (
  if fst a < fst b then
    (le-spvec(snd a,[]) & le-spmat(as, b#bs))
  else (if (fst b < fst a) then
    (le-spvec([], snd b) & le-spmat(a#as, bs))
  else
    (le-spvec(snd a, snd b) & le-spmat (as, bs))))

constdefs
disj-matrices :: ('a::zero) matrix => 'a matrix => bool
disj-matrices A B == (! j i. (Rep-matrix A j i ≠ 0) → (Rep-matrix B j i = 0)) & (! j i. (Rep-matrix B j i ≠ 0) → (Rep-matrix A j i = 0))

⟨ML⟩

lemma disj-matrices-contr1: disj-matrices A B ⇒ Rep-matrix A j i ≠ 0 ⇒ Rep-matrix B j i = 0
⟨proof⟩

lemma disj-matrices-contr2: disj-matrices A B ⇒ Rep-matrix B j i ≠ 0 ⇒ Rep-matrix A j i = 0
⟨proof⟩

lemma disj-matrices-add: disj-matrices A B ⇒ disj-matrices C D ⇒ disj-matrices A D ⇒ disj-matrices B C ⇒
(A + B <= C + D) = (A <= C & B <= (D::('a::lordered-ab-group) matrix))
⟨proof⟩

lemma disj-matrices-zero1[simp]: disj-matrices 0 B
⟨proof⟩

lemma disj-matrices-zero2[simp]: disj-matrices A 0
⟨proof⟩

lemma disj-matrices-commute: disj-matrices A B = disj-matrices B A

```

$\langle \text{proof} \rangle$

lemma *disj-matrices-add-le-zero*: $\text{disj-matrices } A \ B \implies$
 $(A + B \leq 0) = (A \leq 0 \ \& \ (B::('a::\text{ordered-ab-group}) \ \text{matrix}) \leq 0)$
 $\langle \text{proof} \rangle$

lemma *disj-matrices-add-zero-le*: $\text{disj-matrices } A \ B \implies$
 $(0 \leq A + B) = (0 \leq A \ \& \ 0 \leq (B::('a::\text{ordered-ab-group}) \ \text{matrix}))$
 $\langle \text{proof} \rangle$

lemma *disj-matrices-add-x-le*: $\text{disj-matrices } A \ B \implies \text{disj-matrices } B \ C \implies$
 $(A \leq B + C) = (A \leq C \ \& \ 0 \leq (B::('a::\text{ordered-ab-group}) \ \text{matrix}))$
 $\langle \text{proof} \rangle$

lemma *disj-matrices-add-le-x*: $\text{disj-matrices } A \ B \implies \text{disj-matrices } B \ C \implies$
 $(B + A \leq C) = (A \leq C \ \& \ (B::('a::\text{ordered-ab-group}) \ \text{matrix}) \leq 0)$
 $\langle \text{proof} \rangle$

lemma *disj-sparse-row-singleton*: $i \leq j \implies \text{sorted-spvec}((j,y)\#v) \implies \text{disj-matrices}$
 $(\text{sparse-row-vector } v) \ (\text{singleton-matrix } 0 \ i \ x)$
 $\langle \text{proof} \rangle$

lemma *disj-matrices-x-add*: $\text{disj-matrices } A \ B \implies \text{disj-matrices } A \ C \implies \text{disj-matrices}$
 $(A::('a::\text{ordered-ab-group}) \ \text{matrix}) \ (B+C)$
 $\langle \text{proof} \rangle$

lemma *disj-matrices-add-x*: $\text{disj-matrices } A \ B \implies \text{disj-matrices } A \ C \implies \text{disj-matrices}$
 $(B+C) \ (A::('a::\text{ordered-ab-group}) \ \text{matrix})$
 $\langle \text{proof} \rangle$

lemma *disj-singleton-matrices[simp]*: $\text{disj-matrices } (\text{singleton-matrix } j \ i \ x) \ (\text{singleton-matrix}$
 $u \ v \ y) = (j \neq u \mid i \neq v \mid x = 0 \mid y = 0)$
 $\langle \text{proof} \rangle$

lemma *disj-move-sparse-vec-mat[simplified disj-matrices-commute]*:
 $j \leq a \implies \text{sorted-spvec}((a,c)\#as) \implies \text{disj-matrices } (\text{move-matrix } (\text{sparse-row-vector}$
 $b) \ (\text{int } j) \ i) \ (\text{sparse-row-matrix } as)$
 $\langle \text{proof} \rangle$

lemma *disj-move-sparse-row-vector-twice*:
 $j \neq u \implies \text{disj-matrices } (\text{move-matrix } (\text{sparse-row-vector } a) \ j \ i) \ (\text{move-matrix}$
 $(\text{sparse-row-vector } b) \ u \ v)$
 $\langle \text{proof} \rangle$

lemma *le-spvec-iff-sparse-row-le[rule-format]*: $(\text{sorted-spvec } a) \longrightarrow (\text{sorted-spvec}$
 $b) \longrightarrow (\text{le-spvec } (a,b)) = (\text{sparse-row-vector } a \leq \text{sparse-row-vector } b)$
 $\langle \text{proof} \rangle$

lemma *le-spvec-empty2-sparse-row[rule-format]*: $(\text{sorted-spvec } b) \longrightarrow (\text{le-spvec } (b,[]))$

$= (\text{sparse-row-vector } b \leq 0))$
 $\langle \text{proof} \rangle$

lemma *le-spmat-empty1-sparse-row*[rule-format]: $(\text{sorted-spmat } b) \longrightarrow (\text{le-spmat } ([], b))$
 $= (0 \leq \text{sparse-row-vector } b))$
 $\langle \text{proof} \rangle$

lemma *le-spmat-iff-sparse-row-le*[rule-format]: $(\text{sorted-spmat } A) \longrightarrow (\text{sorted-spmat } A) \longrightarrow (\text{sorted-spmat } B) \longrightarrow$
 $\text{le-spmat}(A, B) = (\text{sparse-row-matrix } A \leq \text{sparse-row-matrix } B)$
 $\langle \text{proof} \rangle$

$\langle ML \rangle$

consts

$\text{abs-spmat} :: ('a::\text{ordered-ring}) \text{ spmat} \Rightarrow 'a \text{ spmat}$
 $\text{minus-spmat} :: ('a::\text{ordered-ring}) \text{ spmat} \Rightarrow 'a \text{ spmat}$

primrec

$\text{abs-spmat } [] = []$
 $\text{abs-spmat } (a \# as) = (\text{fst } a, \text{abs-spmat } (\text{snd } a)) \# (\text{abs-spmat } as)$

primrec

$\text{minus-spmat } [] = []$
 $\text{minus-spmat } (a \# as) = (\text{fst } a, \text{minus-spmat } (\text{snd } a)) \# (\text{minus-spmat } as)$

lemma *sparse-row-matrix-minus*:

$\text{sparse-row-matrix } (\text{minus-spmat } A) = - (\text{sparse-row-matrix } A)$
 $\langle \text{proof} \rangle$

lemma *Rep-sparse-row-vector-zero*: $x \neq 0 \implies \text{Rep-matrix } (\text{sparse-row-vector } v)$
 $x \cdot y = 0$
 $\langle \text{proof} \rangle$

lemma *sparse-row-matrix-abs*:

$\text{sorted-spmat } A \implies \text{sorted-spmat } A \implies \text{sparse-row-matrix } (\text{abs-spmat } A) = \text{abs}$
 $(\text{sparse-row-matrix } A)$
 $\langle \text{proof} \rangle$

lemma *sorted-spmat-minus-spmat*: $\text{sorted-spmat } A \implies \text{sorted-spmat } (\text{minus-spmat } A)$
 $\langle \text{proof} \rangle$

lemma *sorted-spmat-abs-spmat*: $\text{sorted-spmat } A \implies \text{sorted-spmat } (\text{abs-spmat } A)$
 $\langle \text{proof} \rangle$

lemma *sorted-spmat-minus-spmat*: $\text{sorted-spmat } A \implies \text{sorted-spmat } (\text{minus-spmat } A)$
 $\langle \text{proof} \rangle$

lemma *sorted-spmat-abs-spmat*: *sorted-spmat A* \implies *sorted-spmat (abs-spmat A)*
 ⟨proof⟩

constdefs

diff-spmat :: ('a::lordered-ring) *spmat* \Rightarrow 'a *spmat* \Rightarrow 'a *spmat*
diff-spmat A B == *add-spmat (A, minus-spmat B)*

lemma *sorted-spmat-diff-spmat*: *sorted-spmat A* \implies *sorted-spmat B* \implies *sorted-spmat*
(diff-spmat A B)
 ⟨proof⟩

lemma *sorted-spvec-diff-spmat*: *sorted-spvec A* \implies *sorted-spvec B* \implies *sorted-spvec*
(diff-spmat A B)
 ⟨proof⟩

lemma *sparse-row-diff-spmat*: *sparse-row-matrix (diff-spmat A B)* = (*sparse-row-matrix*
A) - (*sparse-row-matrix B*)
 ⟨proof⟩

constdefs

sorted-sparse-matrix :: 'a *spmat* \Rightarrow *bool*
sorted-sparse-matrix A == (*sorted-spvec A*) & (*sorted-spmat A*)

lemma *sorted-sparse-matrix-imp-spvec*: *sorted-sparse-matrix A* \implies *sorted-spvec A*
 ⟨proof⟩

lemma *sorted-sparse-matrix-imp-spmat*: *sorted-sparse-matrix A* \implies *sorted-spmat*
A
 ⟨proof⟩

lemmas *sorted-sp-simps* =
sorted-spvec.simps
sorted-spmat.simps
sorted-sparse-matrix-def

lemma *bool1*: (\neg *True*) = *False* ⟨proof⟩

lemma *bool2*: (\neg *False*) = *True* ⟨proof⟩

lemma *bool3*: ((*P*::*bool*) \wedge *True*) = *P* ⟨proof⟩

lemma *bool4*: (*True* \wedge (*P*::*bool*)) = *P* ⟨proof⟩

lemma *bool5*: ((*P*::*bool*) \wedge *False*) = *False* ⟨proof⟩

lemma *bool6*: (*False* \wedge (*P*::*bool*)) = *False* ⟨proof⟩

lemma *bool7*: ((*P*::*bool*) \vee *True*) = *True* ⟨proof⟩

lemma *bool8*: (*True* \vee (*P*::*bool*)) = *True* ⟨proof⟩

lemma *bool9*: ((*P*::*bool*) \vee *False*) = *P* ⟨proof⟩

lemma *bool10*: (*False* \vee (*P*::*bool*)) = *P* ⟨proof⟩

lemmas *boolarith* = *bool1 bool2 bool3 bool4 bool5 bool6 bool7 bool8 bool9 bool10*

lemma *if-case-eq*: (*if b then x else y*) = (*case b of True => x | False => y*)

$\langle proof \rangle$

consts

$pprt\text{-}spvec :: ('a::\{ordered\text{-}ab\text{-}group\})\ spvec \Rightarrow 'a\ spvec$
 $nprr\text{-}spvec :: ('a::\{ordered\text{-}ab\text{-}group\})\ spvec \Rightarrow 'a\ spvec$
 $pprt\text{-}spmat :: ('a::\{ordered\text{-}ab\text{-}group\})\ spmat \Rightarrow 'a\ spmat$
 $nprr\text{-}spmat :: ('a::\{ordered\text{-}ab\text{-}group\})\ spmat \Rightarrow 'a\ spmat$

primrec

$pprt\text{-}spvec\ [] = []$
 $pprt\text{-}spvec\ (a\#as) = (fst\ a,\ pprt\ (snd\ a))\ \# (pprt\text{-}spvec\ as)$

primrec

$nprr\text{-}spvec\ [] = []$
 $nprr\text{-}spvec\ (a\#as) = (fst\ a,\ nprr\ (snd\ a))\ \# (nprr\text{-}spvec\ as)$

primrec

$pprt\text{-}spmat\ [] = []$
 $pprt\text{-}spmat\ (a\#as) = (fst\ a,\ pprt\text{-}spvec\ (snd\ a))\ \# (pprt\text{-}spmat\ as)$

primrec

$nprr\text{-}spmat\ [] = []$
 $nprr\text{-}spmat\ (a\#as) = (fst\ a,\ nprr\text{-}spvec\ (snd\ a))\ \# (nprr\text{-}spmat\ as)$

lemma $pprt\text{-}add$: $disj\text{-}matrices\ A\ (B::(\text{::}ordered\text{-}ring)\ matrix) \Longrightarrow pprt\ (A+B)$
 $= pprt\ A + pprt\ B$
 $\langle proof \rangle$

lemma $nprr\text{-}add$: $disj\text{-}matrices\ A\ (B::(\text{::}ordered\text{-}ring)\ matrix) \Longrightarrow nprr\ (A+B)$
 $= nprr\ A + nprr\ B$
 $\langle proof \rangle$

lemma $pprt\text{-}singleton[simp]$: $pprt\ (singleton\text{-}matrix\ j\ i\ (x::\text{::}ordered\text{-}ring)) = singleton\text{-}matrix\ j\ i\ (pprt\ x)$
 $\langle proof \rangle$

lemma $nprr\text{-}singleton[simp]$: $nprr\ (singleton\text{-}matrix\ j\ i\ (x::\text{::}ordered\text{-}ring)) = singleton\text{-}matrix\ j\ i\ (nprr\ x)$
 $\langle proof \rangle$

lemma $less\text{-}imp\text{-}le$: $a < b \Longrightarrow a \leq (b::\text{::}order)\ \langle proof \rangle$

lemma $sparse\text{-}row\text{-}vector\text{-}pprt$: $sorted\text{-}spvec\ v \Longrightarrow sparse\text{-}row\text{-}vector\ (pprt\text{-}spvec\ v) = pprt\ (sparse\text{-}row\text{-}vector\ v)$
 $\langle proof \rangle$

lemma *sparse-row-vector-nprt*: $\text{sorted-spvec } v \implies \text{sparse-row-vector } (\text{nprt-spvec } v) = \text{nprt } (\text{sparse-row-vector } v)$
 ⟨proof⟩

lemma *pprt-move-matrix*: $\text{pprt } (\text{move-matrix } (A::('a::\text{lordered-ring}) \text{ matrix}) j i) = \text{move-matrix } (\text{pprt } A) j i$
 ⟨proof⟩

lemma *nprt-move-matrix*: $\text{nprt } (\text{move-matrix } (A::('a::\text{lordered-ring}) \text{ matrix}) j i) = \text{move-matrix } (\text{nprt } A) j i$
 ⟨proof⟩

lemma *sparse-row-matrix-pprt*: $\text{sorted-spvec } m \implies \text{sorted-spmat } m \implies \text{sparse-row-matrix } (\text{pprt-spmat } m) = \text{pprt } (\text{sparse-row-matrix } m)$
 ⟨proof⟩

lemma *sparse-row-matrix-nprt*: $\text{sorted-spvec } m \implies \text{sorted-spmat } m \implies \text{sparse-row-matrix } (\text{nprt-spmat } m) = \text{nprt } (\text{sparse-row-matrix } m)$
 ⟨proof⟩

lemma *sorted-pprt-spvec*: $\text{sorted-spvec } v \implies \text{sorted-spvec } (\text{pprt-spvec } v)$
 ⟨proof⟩

lemma *sorted-nprt-spvec*: $\text{sorted-spvec } v \implies \text{sorted-spvec } (\text{nprt-spvec } v)$
 ⟨proof⟩

lemma *sorted-spvec-pprt-spmat*: $\text{sorted-spvec } m \implies \text{sorted-spvec } (\text{pprt-spmat } m)$
 ⟨proof⟩

lemma *sorted-spvec-nprt-spmat*: $\text{sorted-spvec } m \implies \text{sorted-spvec } (\text{nprt-spmat } m)$
 ⟨proof⟩

lemma *sorted-spmat-pprt-spmat*: $\text{sorted-spmat } m \implies \text{sorted-spmat } (\text{pprt-spmat } m)$
 ⟨proof⟩

lemma *sorted-spmat-nprt-spmat*: $\text{sorted-spmat } m \implies \text{sorted-spmat } (\text{nprt-spmat } m)$
 ⟨proof⟩

constdefs

$\text{mult-est-spmat} :: ('a::\text{lordered-ring}) \text{ spmat} \Rightarrow 'a \text{ spmat} \Rightarrow 'a \text{ spmat} \Rightarrow 'a \text{ spmat} \Rightarrow 'a \text{ spmat}$
 $\text{mult-est-spmat } r1 \ r2 \ s1 \ s2 ==$
 $\text{add-spmat } (\text{mult-spmat } (\text{pprt-spmat } s2) (\text{pprt-spmat } r2), \text{add-spmat } (\text{mult-spmat } (\text{pprt-spmat } s1) (\text{nprt-spmat } r2),$
 $\text{add-spmat } (\text{mult-spmat } (\text{nprt-spmat } s2) (\text{pprt-spmat } r1), \text{mult-spmat } (\text{nprt-spmat } s1) (\text{nprt-spmat } r1))))$

lemmas *sparse-row-matrix-op-simps* =
sorted-sparse-matrix-imp-spmat sorted-sparse-matrix-imp-spvec
sparse-row-add-spmat sorted-spvec-add-spmat sorted-spmat-add-spmat
sparse-row-diff-spmat sorted-spvec-diff-spmat sorted-spmat-diff-spmat
sparse-row-matrix-minus sorted-spvec-minus-spmat sorted-spmat-minus-spmat
sparse-row-mult-spmat sorted-spvec-mult-spmat sorted-spmat-mult-spmat
sparse-row-matrix-abs sorted-spvec-abs-spmat sorted-spmat-abs-spmat
le-spmat-iff-sparse-row-le
sparse-row-matrix-pprt sorted-spvec-pprt-spmat sorted-spmat-pprt-spmat
sparse-row-matrix-nprt sorted-spvec-nprt-spmat sorted-spmat-nprt-spmat

lemma *zero-eq-Numeral0*: $(0::\text{number-ring}) = \text{Numeral0}$ *<proof>*

lemmas *sparse-row-matrix-arith-simps*[*simplified zero-eq-Numeral0*] =
mult-spmat.simps mult-spvec-spmat.simps
addmult-spvec.simps
smult-spvec-empty smult-spvec-cons
add-spmat.simps add-spvec.simps
minus-spmat.simps minus-spvec.simps
abs-spmat.simps abs-spvec.simps
diff-spmat-def
le-spmat.simps le-spvec.simps
pprt-spmat.simps pprt-spvec.simps
nprt-spmat.simps nprt-spvec.simps
mult-est-spmat-def

lemma *spm-mult-le-dual-prts*:

assumes

sorted-sparse-matrix A1
sorted-sparse-matrix A2
sorted-sparse-matrix c1
sorted-sparse-matrix c2
sorted-sparse-matrix y
sorted-sparse-matrix r1
sorted-sparse-matrix r2
sorted-spvec b
le-spmat ([], y)
sparse-row-matrix A1 ≤ A
A ≤ sparse-row-matrix A2
sparse-row-matrix c1 ≤ c
c ≤ sparse-row-matrix c2
sparse-row-matrix r1 ≤ x
x ≤ sparse-row-matrix r2
*A * x ≤ sparse-row-matrix (b::('a::lordered-ring) spmat)*

shows

```

  c * x ≤ sparse-row-matrix (add-spmat (mult-spmat y b,
    (let s1 = diff-spmat c1 (mult-spmat y A2); s2 = diff-spmat c2 (mult-spmat y
A1) in
      add-spmat (mult-spmat (pprt-spmat s2) (pprt-spmat r2), add-spmat (mult-spmat
(pprt-spmat s1) (nprrt-spmat r2),
      add-spmat (mult-spmat (nprrt-spmat s2) (pprt-spmat r1), mult-spmat (nprrt-spmat
s1) (nprrt-spmat r1))))))
  ⟨proof⟩

```

lemma *spm-mult-le-dual-prts-no-let*:

```

assumes
  sorted-sparse-matrix A1
  sorted-sparse-matrix A2
  sorted-sparse-matrix c1
  sorted-sparse-matrix c2
  sorted-sparse-matrix y
  sorted-sparse-matrix r1
  sorted-sparse-matrix r2
  sorted-spmat b
  le-spmat ([], y)
  sparse-row-matrix A1 ≤ A
  A ≤ sparse-row-matrix A2
  sparse-row-matrix c1 ≤ c
  c ≤ sparse-row-matrix c2
  sparse-row-matrix r1 ≤ x
  x ≤ sparse-row-matrix r2
  A * x ≤ sparse-row-matrix (b::('a::lordered-ring) spmat)
shows
  c * x ≤ sparse-row-matrix (add-spmat (mult-spmat y b,
    mult-est-spmat r1 r2 (diff-spmat c1 (mult-spmat y A2)) (diff-spmat c2 (mult-spmat
y A1))))
  ⟨proof⟩

```

end

theory *FloatSparseMatrix* **imports** *Float SparseMatrix* **begin**

end

```

theory Cplex
imports FloatSparseMatrix
uses Cplex-tools.ML CplexMatrixConverter.ML FloatSparseMatrixBuilder.ML fspmlp.ML
begin

```

end

```

theory MatrixLP
imports Cplex
begin

constdefs
  list-case-compute :: 'b list  $\Rightarrow$  'a  $\Rightarrow$  ('b  $\Rightarrow$  'b list  $\Rightarrow$  'a)  $\Rightarrow$  'a
  list-case-compute l a f == list-case a f l

lemma list-case-compute: list-case = ( $\lambda$  (a::'a) f (l::'b list). list-case-compute l a
f)
  <proof>

lemma list-case-compute-empty: list-case-compute ([]::'b list) = ( $\lambda$  (a::'a) f. a)
  <proof>

lemma list-case-compute-cons: list-case-compute (u#v) = ( $\lambda$  (a::'a) f. (f (u::'b)
v))
  <proof>

lemma If-True: (If True) = ( $\lambda$  x y. x)
  <proof>

lemma If-False: (If False) = ( $\lambda$  x y. y)
  <proof>

lemma Let-compute: Let (x::'a) f = ((f x)::'b)
  <proof>

lemma fst-compute: fst (a::'a, b::'b) = a
  <proof>

lemma snd-compute: snd (a::'a, b::'b) = b
  <proof>

lemma bool1: ( $\neg$  True) = False <proof>
lemma bool2: ( $\neg$  False) = True <proof>
lemma bool3: ((P::bool)  $\wedge$  True) = P <proof>
lemma bool4: (True  $\wedge$  (P::bool)) = P <proof>
lemma bool5: ((P::bool)  $\wedge$  False) = False <proof>
lemma bool6: (False  $\wedge$  (P::bool)) = False <proof>
lemma bool7: ((P::bool)  $\vee$  True) = True <proof>
lemma bool8: (True  $\vee$  (P::bool)) = True <proof>
lemma bool9: ((P::bool)  $\vee$  False) = P <proof>

```

```

lemma bool10: (False  $\vee$  (P::bool)) = P <proof>
lemmas boolarith = bool1 bool2 bool3 bool4 bool5 bool6 bool7 bool8 bool9 bool10

lemmas float-arith = Float.arith
lemmas sparse-row-matrix-arith-simps = SparseMatrix.sparse-row-matrix-arith-simps
lemmas sorted-sp-simps = SparseMatrix.sorted-sp-simps
lemmas fst-snd-conv = Product-Type.fst-conv Product-Type.snd-conv

end

```